## From the ax + b Lie Algebra to the Alexander Polynomial and Beyond Dror Bar–Natan, Chicago, September 2010 http://www.math.toronto.edu/~drorbn/Talks/Chicago–1009/

Dror Bar–Natan, Chicago, September 2010	http://www.math.toronto.edu/~drorbn/Talks/Chicago-1009/
Abstract. I will present the simplest-ever "quantum" formula	An Alexander Reminder.
for the Alexander polynomial, using only the unique two di-	Number the arrows $1, \ldots, n, $
mensional non-commutative Lie algebra (the one associated	let $t_i, h_j$ be the tail and head $(a_1   (a_2   (a_3   (a_4))))$
with the " $ax + b$ " Lie group). After introducing the "Euler	
technique" and some diagrammatic calculus I will sketch the	of arrow j, and let $s_j \in \pm 1$ be $\begin{pmatrix} 1-X+(-1) & X & 0 & 0 \\ 1-X^{-1} & 0 & -1 & X^{-1} \end{pmatrix}$
proof of the said formula, and following that, I will present a	arcs $a_{\alpha}$ by arrow heads, and $\begin{pmatrix} 1 & 1 \\ 0 & -1 & X & 1-X \end{pmatrix}$
long list of extensions, generalizations, and dreams.	let $\alpha(p)$ be "the arc of point $p$ ". Let $R \in M_{n \times (n+1)}$ be the
The 2D Lie Algebra. Let $\mathfrak{g} = \mathfrak{lie}(x^1, x^2)/[x^1, x^2] = x^2$ , let	
$\mathfrak{g}^* = \langle \phi_1, \phi_2 \rangle$ with $\phi_i(x^j) = \delta_i^j$ , let $I\mathfrak{g} = \mathfrak{g}^* \rtimes \mathfrak{g}$ so	matrix whose j'th row has $-1$ in column $\alpha(h_j)$ and $1 - X^{s_j}$
$[\phi_i, \phi_i] = [\phi_1, x^i] = 0$ while $[x^1, \phi_2] = -\phi_2$ and $[x^2, \phi_2] = \phi_1$ .	in column $\alpha(t_j)$ and $X^{s_j}$ in column $\alpha(h_j) + 1$ , and let M be
Let $r = Id = \phi_1 \otimes x^1 + \phi_2 \otimes x^2 \in \mathfrak{g}^* \otimes \mathfrak{g} \subset I\mathfrak{g} \otimes I\mathfrak{g}.$	$R$ with a column removed. Then $A(X) = \det(M)$ .
Let $\mathcal{U} = \{ \text{words in } I\mathfrak{a} \} / ab - ba = [a, b] \text{ degree-completed}$	An Euler Interlude. If you know brackets, how do you test
with respect to deg $\phi_i = 1$ and deg $x^i = 0$ (so $\mathcal{U} \equiv$	exponentials? When's $e^{A}e^{B} = e^{C}e^{D}$ ?
(power series is 4 variables)). Let $R = \exp(r) \in \mathcal{U} \otimes \mathcal{U}$ .	Bad Idea. Take log and use BCH. You'll want to cry.
	Clever Idea. Let $E$ be the Euler derivation, which mul-
The Invariant. Define $Z$ : $+$ $ +$	tiplies each element by its degree (e.g. on $\mathbb{Q}[\![\phi]\!]$ , $Ef =$
$\{\text{long knots}\} \rightarrow \mathcal{U} \text{ by mapping}$	$\phi \partial_{\phi} f$ , so $Ee^{\phi} = \phi e^{\phi}$ ). Apply $\tilde{E}\zeta := \zeta^{-1}E\zeta$ : $\tilde{E}(e^{A}e^{B}) =$
every $\pm$ -crossing to $R^{\pm 1}$ :	$e^{-B}e^{-A}(e^{A}Ae^{B} + e^{A}e^{B}B) = e^{-B}Ae^{B} + B = e^{-\operatorname{ad} B}(A) + B.$
	e e (e Ae + e e D) - e Ae + D - e (A) + D.
	"Uninterpreting" Diagrams. Make $Z^w : \mathcal{K}^w \to \mathcal{A}^w \to \mathcal{U}$ , with
Alexander	
	$($ $($ $)$ $($ $)$ $($ $2$ in 1 out vertices $)$ $($ $\overrightarrow{STU}, \overrightarrow{AS},$
$\cdots + \frac{1}{2!} \frac{(-1)^3}{3!} \frac{1}{1!} (\phi_2 \phi_1) (\phi_2 \phi_1 \phi_2) (x^2 x^1) (x^1) (x^2 x^1 x^2) (\phi_1) + \cdots$	$\mathcal{A}^w = \left( \begin{array}{c} \mathcal{A}^w = \left( \begin{array}{c} \mathcal{A}^w \right) \\ \mathcal{A}^w = \left( \begin{array}{c} \mathcal{A}^w = \left( \begin{array}{c} \mathcal{A}^w \right) \\ \mathcal{A}^w = \left( \begin{array}{c} \mathcal{A}^w = \left( \begin{array}{c} \mathcal{A}^w \right) \\ \mathcal{A}^w = \left( \begin{array}{c} \mathcal{A}^w = \left( \begin{array}{c} \mathcal{A}^w \right) \\ \mathcal{A}^w = \left( \begin{array}{c} \mathcal{A}^w = \left( \right) $
Near Theorem Zicinversiont and it is acceptially the Aleven	relations
Near Theorem. Z is invariant, and it is essentially the Alexan-	
der polynomial; with $N = \exp(\overline{l} \phi_i x^i + \overline{l} x^i \phi_i) =: \exp(SL)$ ,	
$Z(K) = N \cdot \left(A(K)(e^{\phi_1})\right)^{-1} \tag{1}$	$\overrightarrow{STU}_1$ : = $\overrightarrow{TU}_2$ : = $\overrightarrow{TU}_2$ :
Invariance. "The identity is an invariant tensor":	
invariance. The identity is an invariant tensor .	
	$\overrightarrow{STU}_3 = \text{TC: } 0 = \bigcirc - \bigcirc \overrightarrow{IHX}$
$\begin{array}{c} \hline \\ \hline $	$\mathcal{K}^w = CA \left\langle \begin{array}{c} & \\ \end{array} \right\rangle / R23, OC $
The Euler Prelude. Apply $\tilde{E}\zeta := \zeta^{-1}E\zeta$ to (1):	$= PA \left\langle \checkmark \right\rangle / R23, VR123, D, OC $
$+Z^{-1}$ $-Z^{-1}$ $($	X     X   X     X   X     X   X     X   X     X   X   X     X
$ \begin{array}{c c} & & & \\ \hline \\ \hline$	
<i>T i T i T i T</i>	
$\stackrel{?}{=} SL$	R3 VR3 D OC
$+Z^{-1} \left( \int -\phi_1 \operatorname{tr} \left( M^{-1} \frac{d}{d\phi_1} M(e^{\phi_1}) \right) \right)$	$Z^w$ is a UFTI on w-knots! It extends to links and tangles,
$x^i \phi_i$	is well behaved under compositions and cables, and remains
Some Relations. $\phi_i x^i$ , $x^i \phi_i$ , $\phi_1$ are central, $x^i \phi_i - \phi_i x^i = \phi_1$ ,	computable for tangles. It contains Burau, Gassner, and
$\begin{bmatrix} x^j, \phi_i \end{bmatrix} = \delta_i^j \phi_1 - \delta_1^j \phi_i \text{ or } j  i  i$	Cimasoni-Turaev in natural ways, and it contains the MVA
	though my understanding of the latter is incomplete.
so $ -  =  =  -  -  $	w-Knots.
S et a the state of the state o	
$ \underbrace{ \begin{array}{c} \\ \\ \end{array}}^{s} \underbrace{ \\ \\ \end{array}}^{s} - e^{s\phi_1} \underbrace{ \begin{array}{c} \\ \\ \end{array}}^{s} \underbrace{ \\ \end{array}}^{s} \underbrace{ \\ \end{array}}^{s} = (1 - e^{s\phi_1}) \underbrace{ \begin{array}{c} \\ \\ \end{array}}^{s} \underbrace{ \\ \end{array}}^{s} \underbrace{ \\ \end{array}}^{s} = \phi_1 $	2D Symbol
and the tamed "tails commute" (TC): $\sqrt{S} \times \sqrt{S} \sqrt{S}$	
and the famed "tails commute" (TC): $\sqrt{s} = \sqrt{s}$	
	Crossing
Near Proof. Let $\lambda_{\alpha j}$ be a red ar- $\lambda_{12}$	Crossing Wen W Vertices
Near Proof. Let $\lambda_{\alpha j}$ be a red arrow with tail at $a_{\alpha}$ and head just $1 -2 -3$	Crossing Virtual crossing Movie
Near Proof. Let $\lambda_{\alpha j}$ be a red arrow with tail at $a_{\alpha}$ and head just left of $h_j$ . Let $\Lambda = (\lambda_{\alpha j})$ . Then	Crossing Virtual crossing Movie
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Near Proof. Let $\lambda_{\alpha j}$ be a red arrow with tail at $a_{\alpha}$ and head just left of $h_j$ . Let $\Lambda = (\lambda_{\alpha j})$ . Then roughly $R\Lambda = \phi_1 I$ so roughly, $\Lambda = R^{-1}\phi_1$ . The rest is book-keeping that I haven't finished	Crossing Virtual crossing Movie
Near Proof. Let $\lambda_{\alpha j}$ be a red arrow with tail at $a_{\alpha}$ and head just left of $h_j$ . Let $\Lambda = (\lambda_{\alpha j})$ . Then roughly $R\Lambda = \phi_1 I$ so roughly, $(a_1 ) (a_2) (a_3) (a_4)$ $\Lambda = R^{-1}\phi_1$ . The rest is book-keeping that I haven't finished yet, yet with which my computer agrees fully.	Crossing Wen W Cap Wen W C
Near Proof. Let $\lambda_{\alpha j}$ be a red arrow with tail at $a_{\alpha}$ and head just left of $h_j$ . Let $\Lambda = (\lambda_{\alpha j})$ . Then roughly $R\Lambda = \phi_1 I$ so roughly, $\Lambda = R^{-1}\phi_1$ . The rest is book-keeping that I haven't finished	Crossing Virtual crossing Movie Crossing Wen Wen Vertices smooth singular There's 1D in 4D, non-trivial given 2D, and there are ops Dream. $Z^w$ extends to virtual knots as $Z^v : \mathcal{K}^v \to \mathcal{A}^v$ , with
Near Proof. Let $\lambda_{\alpha j}$ be a red arrow with tail at $a_{\alpha}$ and head just left of $h_j$ . Let $\Lambda = (\lambda_{\alpha j})$ . Then roughly $R\Lambda = \phi_1 I$ so roughly, $\Lambda = R^{-1}\phi_1$ . The rest is book-keeping that I haven't finished yet, yet with which my computer agrees fully. I don't understand the Alexander polynomial! "God created the knots, all else in	Crossing Wen W Cap Wen W C
Near Proof. Let $\lambda_{\alpha j}$ be a red arrow with tail at $a_{\alpha}$ and head just left of $h_j$ . Let $\Lambda = (\lambda_{\alpha j})$ . Then roughly $R\Lambda = \phi_1 I$ so roughly, $\begin{pmatrix} 1 & -2 & \lambda_{12} \\ 1 & -2 & \lambda_{3} \\ a_1 & (a_2)(a_3)(a_4) \end{pmatrix}$ $\Lambda = R^{-1}\phi_1$ . The rest is book-keeping that I haven't finished yet, yet with which my computer agrees fully. I don't understand the Alexander polynomial!	Crossing Virtual crossing Movie Crossing Wen Wen Vertices smooth singular There's 1D in 4D, non-trivial given 2D, and there are ops Dream. $Z^w$ extends to virtual knots as $Z^v : \mathcal{K}^v \to \mathcal{A}^v$ , with