

# Meta-Groups, Meta-Bicrossed-Products, and the Alexander Polynomial, 1

Dror Bar-Natan in Montreal, June 2013.

<http://www.math.toronto.edu/~drorbn/Talks/Montreal-1306/>



**Abstract.** I will define “meta-groups” and explain how one specific meta-group, which in itself is a “meta-bicrossed-product”, gives rise to an “ultimate Alexander invariant” of tangles, that contains the Alexander polynomial (multivariable, if you wish), has extremely good composition properties, is evaluated in a topologically meaningful way, and is least-wasteful in a computational sense. If you believe in categorification, that’s a wonderful playground.

This work is closely related to work by Le Dimet (Comment. Math. Helv. **67** (1992) 306–315), Kirk, Livingston and Wang (arXiv:math/9806035) and Cimasoni and Turaev (arXiv:math.GT/0406269).

See also Dror Bar-Natan and Sam Selmani, *Meta-Monoids, Meta-Bicrossed Products, and the Alexander Polynomial*, arXiv:1302.5689.

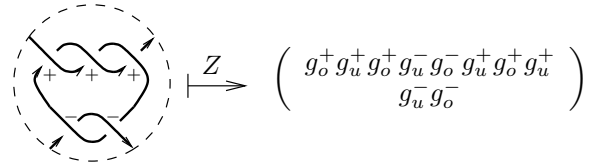
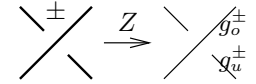


Sam Selmani

## Alexander Issues.

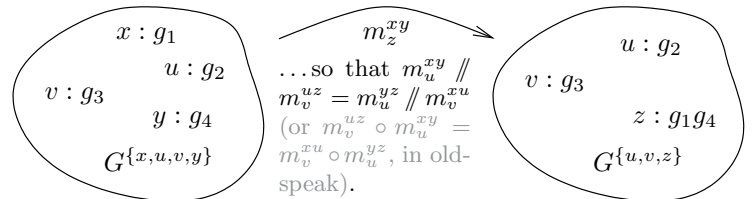
- Quick to compute, but computation departs from topology.
- Extends to tangles, but at an exponential cost.
- Hard to categorify.

**Idea.** Given a group  $G$  and two “YB” pairs  $R^\pm = (g_o^\pm, g_u^\pm) \in G^2$ , map them to xings and “multiply along”, so that



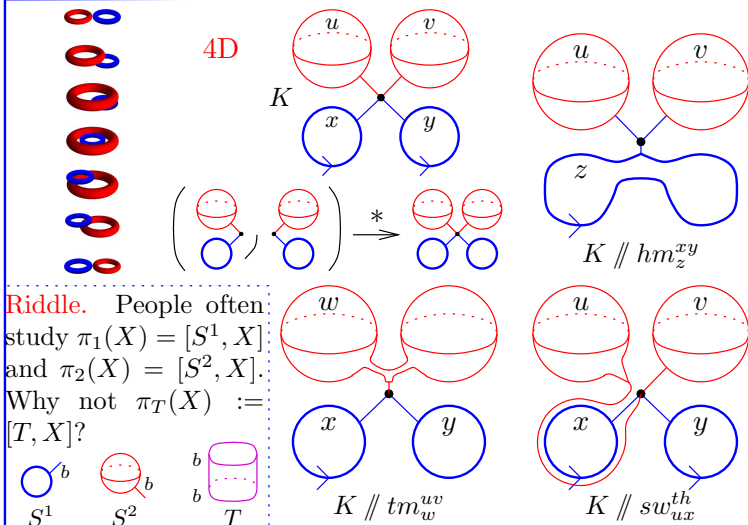
**This Fails!** R2 implies that  $g_o^\pm g_o^\mp = e = g_u^\pm g_u^\mp$  and then R3 implies that  $g_o^\pm$  and  $g_u^\pm$  commute, so the result is a simple counting invariant.

**A Group Computer.** Given  $G$ , can store group elements and perform operations on them:

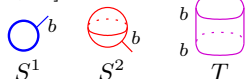


Also has  $S_x$  for inversion,  $e_x$  for unit insertion,  $d_x$  for register deletion,  $\Delta_{xy}^z$  for element cloning,  $\rho_y^x$  for renamings, and  $(D_1, D_2) \mapsto D_1 \cup D_2$  for merging, and many obvious composition axioms relating those.

$$P = \{x : g_1, y : g_2\} \Rightarrow P = \{d_y P\} \cup \{d_x P\}$$



**Riddle.** People often study  $\pi_1(X) = [S^1, X]$  and  $\pi_2(X) = [S^2, X]$ . Why not  $\pi_T(X) := [T, X]$ ?



**A Meta-Group.** Is a similar “computer”, only its internal structure is unknown to us. Namely it is a collection of sets  $\{G_\gamma\}$  indexed by all finite sets  $\gamma$ , and a collection of operations  $m_z^{xy}$ ,  $S_x$ ,  $e_x$ ,  $d_x$ ,  $\Delta_{xy}^z$  (sometimes),  $\rho_y^x$ , and  $\cup$ , satisfying the exact same *linear* properties.

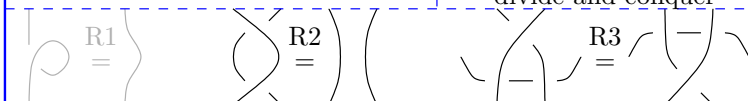
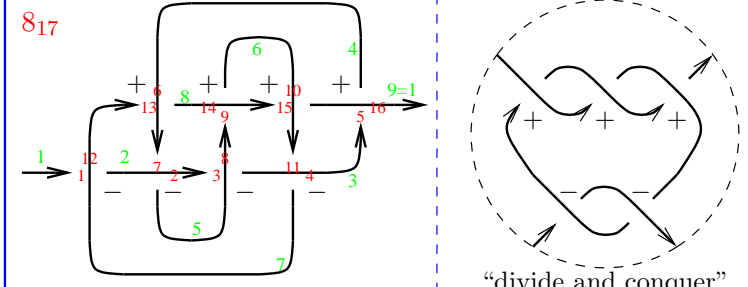
**Example 0.** The non-meta example,  $G_\gamma := G^\gamma$ .

**Example 1.**  $G_\gamma := M_{\gamma \times \gamma}(\mathbb{Z})$ , with simultaneous row and column operations, and “block diagonal” merges. Here if

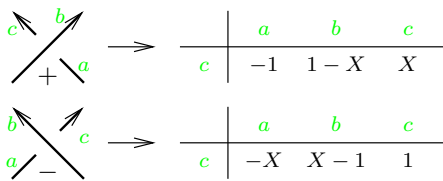
$P = \begin{pmatrix} x & a & b \\ y & c & d \end{pmatrix}$  then  $d_y P = (x : a)$  and  $d_x P = (y : d)$  so

$\{d_y P\} \cup \{d_x P\} = \begin{pmatrix} x & a & 0 \\ y & 0 & d \end{pmatrix} \neq P$ . So this  $G$  is truly meta.

**Claim.** From a meta-group  $G$  and YB elements  $R^\pm \in G_2$  we can construct a knot/tangle invariant.



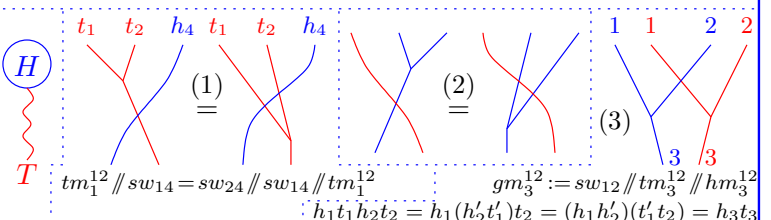
**A Standard Alexander Formula.** Label the arcs  $1$  through  $(n + 1) = 1$ , make an  $n \times n$  matrix as below, delete one row and one column, and compute the determinant:



$$\begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & x-1 & 0 & -x \\ -1 & x & 0 & 0 & 0 & 0 & 0 & 1-x & 0 \\ 0 & -1 & x & 0 & 1-x & 0 & 0 & 0 & 0 \\ x-1 & 0 & -x & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1-x & 0 & -1 & x & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & -x & 1 & 0 & 0 & x-1 \\ 0 & 0 & 1-x & 0 & 0 & 0 & -1 & x & 0 \\ 0 & 0 & 0 & x-1 & 0 & 0 & -x & 1 & 0 \end{pmatrix} \quad [[1 ;; 7, 1 ;; 7]] \quad // \text{ Det}$$

$$-1 + 4x - 8x^2 + 11x^3 - 8x^4 + 4x^5 - x^6$$

**Bicrossed Products.** If  $G = HT$  is a group presented as a product of two of its subgroups, with  $H \cap T = \{e\}$ , then also  $G = TH$  and  $G$  is determined by  $H$ ,  $T$ , and the “swap” map  $sw^{th} : (t, h) \mapsto (h', t')$  defined by  $th = h't'$ . The map  $sw$  satisfies (1) and (2) below; conversely, if  $sw : T \times H \rightarrow H \times T$  satisfies (1) and (2) (+ lesser conditions), then (3) defines a group structure on  $H \times T$ , the “bicrossed product”.



# Meta-Groups, Meta-Bicrossed-Products, and the Alexander Polynomial, 2

A **Meta-Bicrossed-Product** is a collection of sets  $\beta(\eta, \tau)$  and operations  $tm_w^{uv}$ ,  $hm_z^{xy}$  and  $sw_{ux}^{th}$  (and lesser ones), such that  $tm$  and  $hm$  are “associative” and (1) and (2) hold (+ lesser conditions). A meta-bicrossed-product defines a meta-group with  $G_\gamma := \beta(\gamma, \gamma)$  and  $gm$  as in (3).

**Example.** Take  $\beta(\eta, \tau) = M_{\tau \times \eta}(\mathbb{Z})$  with row operations for the tails, column operations for the heads, and a trivial swap.

**$\beta$  Calculus.** Let  $\beta(\eta, \tau)$  be

$$\left\{ \begin{array}{c|ccc} \omega & h_1 & h_2 & \cdots \\ t_1 & \alpha_{11} & \alpha_{12} & \cdot \\ t_2 & \alpha_{21} & \alpha_{22} & \cdot \\ \vdots & \cdot & \cdot & \cdot \end{array} \middle| \begin{array}{l} h_j \in \eta, t_i \in \tau, \text{ and } \omega \text{ and} \\ \text{the } \alpha_{ij} \text{ are rational func-} \\ \text{tions in a variable } X \text{ with} \\ \omega(1) = 1 \text{ and } \alpha_{ij}(1) = 0 \end{array} \right\},$$

$$tm_w^{uv} : \begin{array}{c|c} \omega & \cdots \\ t_u & \alpha \\ t_v & \beta \\ \vdots & \gamma \end{array} \mapsto \begin{array}{c|c} \omega & \cdots \\ t_w & \alpha + \beta \\ \vdots & \gamma \end{array}, \quad \begin{array}{c|c} \omega_1 & \eta_1 \\ \tau_1 & \alpha_1 \end{array} \cup \begin{array}{c|c} \omega_2 & \eta_2 \\ \tau_2 & \alpha_2 \end{array} = \begin{array}{c|c} \omega_1\omega_2 & \eta_1 \eta_2 \\ \tau_1 & \alpha_1 \ 0 \\ \tau_2 & 0 \ \alpha_2 \end{array},$$

$$hm_z^{xy} : \begin{array}{c|ccc} \omega & h_x & h_y & \cdots \\ \vdots & \alpha & \beta & \gamma \end{array} \mapsto \begin{array}{c|ccc} \omega & h_z & & \cdots \\ \vdots & \alpha + \beta + \langle \alpha \rangle \beta & & \gamma \end{array},$$

$$sw_{ux}^{th} : \begin{array}{c|ccc} \omega & h_x & \cdots \\ t_u & \alpha & \beta \\ \vdots & \gamma & \delta \end{array} \mapsto \begin{array}{c|ccc} \omega \epsilon & h_x & \cdots \\ t_u & \alpha(1 + \langle \gamma \rangle / \epsilon) & \beta(1 + \langle \gamma \rangle / \epsilon) \\ \vdots & \gamma / \epsilon & \delta - \gamma \beta / \epsilon \end{array},$$

where  $\epsilon := 1 + \alpha$  and  $\langle c \rangle := \sum_i c_i$ , and let

$$R_{ab}^p := \begin{array}{c|cc} 1 & h_a & h_b \\ t_a & 0 & X-1 \\ t_b & 0 & 0 \end{array} \quad R_{ab}^m := \begin{array}{c|cc} 1 & h_a & h_b \\ t_a & 0 & X^{-1}-1 \\ t_b & 0 & 0 \end{array}.$$

**Theorem.**  $Z^\beta$  is a tangle invariant (and more). Restricted to knots, the  $\omega$  part is the Alexander polynomial. On braids, it is equivalent to the Burau representation. A variant for links contains the multivariable Alexander polynomial.

**Why Happy?** • Applications to w-knots.

• Everything that I know about the Alexander polynomial can be expressed cleanly in this language (even if without proof), except HF, but including genus, ribboness, cabling, v-knots, knotted graphs, etc., and there’s potential for vast generalizations.

• The least wasteful “Alexander for tangles” I’m aware of.

• Every step along the computation is the invariant of something.

• Fits on one sheet, including implementation & propaganda.



**Further meta-monoids.**  $\Pi$  (and variants),  $\mathcal{A}$  (and quotients),  $vT$ , ...

**Further meta-bicrossed-products.**  $\Pi$  (and variants),  $\vec{\mathcal{A}}$  (and quotients),  $M_0$ ,  $M$ ,  $\mathcal{K}^{bh}$ ,  $\mathcal{K}^{rbh}$ , ...

**Meta-Lie-algebras.**  $\mathcal{A}$  (and quotients),  $\mathcal{S}$ , ...

**Meta-Lie-bialgebras.**  $\vec{\mathcal{A}}$  (and quotients), ...

I don’t understand the relationship between  $gr$  and  $H$ , as it appears, for example, in braid theory.

I mean business!

```

SSimp = Factor; SetAttributes[BCollect, Listable];
BCollect[B[omega_, A_] := B[SSimp[omega],
Collect[A, h, Collect[t, t, SSimp] &]];
BForm[B[omega_, A_] := Module[{ts, hs, M},
ts = Union[Cases[B[omega, A], t_u -> u, Infinity]];
hs = Union[Cases[B[omega, A], h_x -> x, Infinity]];
M = Outer[SSimp[Coefficient[A, h_x t_u]] &, hs, ts];
PrependTo[M, t_u & /@ ts];
M = Prepend[Transpose[M], Prepend[h_x & /@ hs, omega]];
MatrixForm[M]];
BForm[else_] := else /. B_B -> BForm[B];
Format[B_B, StandardForm] := BForm[B];
    
```

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{omega_} := # /. t_u -> 1;
tm_u_v_w[omega_] := BCollect[B /. t_u -> t_u];
hm_x_y_z[B[omega_, A_] := Module[
{alpha = D[A, h_x], beta = D[A, h_y], gamma = A /. h_x -> 0},
B[omega, (alpha + (1 + alpha) beta) h_x + gamma] // BCollect];
sw_u_x[B[omega_, A_] := Module[{alpha, beta, gamma, delta, epsilon},
alpha = Coefficient[A, h_x t_u]; beta = D[A, t_u] /. h_x -> 0;
gamma = D[A, h_x] /. t_u -> 0; delta = A /. h_x t_u -> 0;
epsilon = 1 + alpha;
B[omega + epsilon, alpha (1 + gamma/epsilon) h_x t_u + beta (1 + gamma/epsilon) t_u
+ gamma/epsilon h_x
+ delta - gamma*beta/epsilon
] // BCollect];
gm_u_v_w[omega_] := beta // sw_ab // hm_u_v_w // tm_u_v_w;
B /. B[omega_, A_] B[omega_2, A2_] := B[omega_1*omega_2, A1+A2];
Rp_u_v := B[1, (X-1) t_u h_v];
Rm_u_v := B[1, (X^{-1}-1) t_u h_v];
    
```

$$\{\beta = \mathbf{B}[\omega, \text{Sum}[\alpha_{10+i+j} t_i h_j, \{\mathbf{i}, \{1, 2, 3\}\}, \{\mathbf{j}, \{4, 5\}\}]\},$$

$$(\beta // tm_{12 \rightarrow 1} // sw_{14}) = (\beta // sw_{24} // sw_{14} // tm_{12 \rightarrow 1})$$

$$\left\{ \begin{array}{c|ccc} \omega & h_4 & h_5 \\ t_1 & \alpha_{14} & \alpha_{15} \\ t_2 & \alpha_{24} & \alpha_{25} \\ t_3 & \alpha_{34} & \alpha_{35} \end{array} \right\}, \text{True} \quad \begin{array}{c} \text{Y} \\ \text{X} \end{array} \stackrel{(1)}{=} \begin{array}{c} \text{X} \\ \text{Y} \end{array} \quad \text{Some testing}$$

$$\left\{ \begin{array}{c|ccc} 1 & h_1 & h_2 \\ t_2 & -\frac{-1+X}{X} & 0 \\ t_3 & -\frac{-1+X}{X} & -\frac{-1+X}{X} \end{array} \right\}, \left\{ \begin{array}{c|ccc} 1 & h_1 & h_2 \\ t_2 & -\frac{-1+X}{X} & 0 \\ t_3 & -\frac{-1+X}{X} & -\frac{-1+X}{X} \end{array} \right\}$$

... divide and conquer!

$$\beta = Rm_{12,1} Rm_{27} Rm_{83} Rm_{4,11} Rp_{16,5} Rp_{6,13} Rp_{14,9} Rp_{10,15} \quad 817$$

$$\begin{pmatrix} 1 & h_1 & h_3 & h_5 & h_7 & h_9 & h_{11} & h_{13} & h_{15} \\ t_2 & 0 & 0 & 0 & -\frac{-1+X}{X} & 0 & 0 & 0 & 0 \\ t_4 & 0 & 0 & 0 & 0 & 0 & -\frac{-1+X}{X} & 0 & 0 \\ t_6 & 0 & 0 & 0 & 0 & 0 & 0 & -1+X & 0 \\ t_8 & 0 & -\frac{-1+X}{X} & 0 & 0 & 0 & 0 & 0 & 0 \\ t_{10} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & -1+X \\ t_{12} & -\frac{-1+X}{X} & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ t_{14} & 0 & 0 & 0 & 0 & -1+X & 0 & 0 & 0 \\ t_{16} & 0 & 0 & -1+X & 0 & 0 & 0 & 0 & 0 \end{pmatrix}$$

$$Do[\beta = \beta // gm_{1k \rightarrow 1}, \{k, 2, 10\}]; \beta \quad 817, \text{ cont.}$$

$$\begin{pmatrix} \frac{1}{X} & h_1 & & h_{11} & & h_{13} & & h_{15} \\ t_1 & -\frac{(-1+X)(1+X)}{X} & & -(-1+X)(1-X+X^2) & & (-1+X)(1-X+X^2) & & -1+X \\ t_{12} & -\frac{-1+X}{X} & & 0 & & 0 & & 0 \\ t_{14} & -1+X & & \frac{(-1+X)^2(1-X+X^2)}{X} & & -\frac{(-1+X)^2(1-X+X^2)}{X} & & 0 \\ t_{16} & -\frac{-1+X}{X} & & (-1+X)^2 & & -\frac{(-1+X)^3}{X} & & 0 \end{pmatrix}$$



James Waddell Alexander

$$Do[\beta = \beta // gm_{1k \rightarrow 1}, \{k, 11, 16\}]; \beta$$

$$\left( -\frac{1-4X+8X^2-11X^3+8X^4-4X^5+X^6}{X^3} \right)$$

**A Partial To Do List.** 1. Where does it *more simply* come from?

2. Remove all the denominators.

3. How do determinants arise in this context?

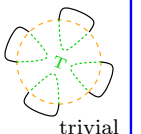
4. Understand links (“meta-conjugacy classes”).

5. Find the “reality condition”.

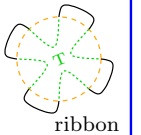
6. Do some “Algebraic Knot Theory”.

7. Categorify.

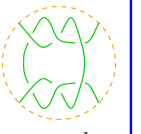
8. Do the same in other natural quotients of the v/w-story.



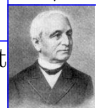
trivial



ribbon



example



“God created the knots, all else in topology is the work of mortals.”  
Leopold Kronecker (modified)

