

```

QZipgs_List@E[L_, Q_, P_] :=

PPQzip@Module[{g, z, zs, c, ys, ηs, qt, zrule, grule, out},
  zs = Table[g*, {g, gs}];
  c = CF[Q /. Alternatives @@ (gs ∪ zs) → 0];
  ys = CF@Table[∂g(Q /. Alternatives @@ zs → 0),
    {g, gs}];
  ηs = CF@Table[∂z(Q /. Alternatives @@ gs → 0), {z, zs}];
  qt = CF@Inverse@Table[Kδz,g* - ∂z,gQ, {g, gs}, {z, zs}];
  zrule = Thread[zs → CF[qt.(zs + ys)]];
  grule = Thread[gs → gs + ηs.qt];
  CF /@ E[L, c + ηs.qt.ys,
  Det[qt] Zipgs[P /. (zrule ∪ grule))]];
];

```

LZip implements the “ L -level zips” on $\mathbb{E}(L, Q, P) = \mathbb{P}\mathbf{e}^{L+Q}$. Such zips regard all of $\mathbb{P}\mathbf{e}^Q$ as a single “ P ”. Here the z ’s are b and α and the g ’s are β and a .

```

LZipgs_List@E[L_, Q_, P_] :=

PPLzip@Module[{g, z, zs, Zs, c, ys, ηs, lt, zrule,
  zrule, grule, Q1, EEQ, EQ},
  zs = Table[g*, {g, gs}];
  Zs = zs /. {b → B, t → T, α → A};
  c = L /. Alternatives @@ (gs ∪ zs) → 0 /.
    Alternatives @@ Zs → 1;
  ys = Table[∂g(L /. Alternatives @@ zs → 0), {g, gs}];
  ηs = Table[∂z(L /. Alternatives @@ gs → 0), {z, zs}];
  lt = Inverse@Table[Kδz,g* - ∂z,gL, {g, gs}, {z, zs}];
  zrule = Thread[zs → lt.(zs + ys)];
  Zrule = Join[zrule,
    zrule /.
    r_Rule :> ((U = r[[1]] /. {b → B, t → T, α → A}) →
      (U /. U2I /. r //.12U));
  grule = Thread[gs → gs + ηs.lt];
  Q1 = Q /. (Zrule ∪ grule);
  EEQ[ps_] :=
    EEQ[ps] =
      PPEEQ@(CF[e-Q1 DThread[{zs, {ps}]] [eQ1] /.
        {Alternatives @@ zs → 0, Alternatives @@ Zs → 1});
  CF@E[c + ηs.lt.ys,
    Q1 /.
      {Alternatives @@ zs → 0, Alternatives @@ Zs → 1},
    Det[lt]
    (Zipgs[({EQ @@ zs) (P /. (Zrule ∪ grule))] /.
      Derivative[ps_][EQ] [__] :> EEQ[ps] /.
      EQ → 1] ];
];

```

```

B0[L_, R_] := L R;
B{is_}[L E, R E] := PPB@Module[{n},
  Times[L /. Table[(v : b | B | t | T | a | x | y)i → vnei,
    {i, {is}}],
  R /. Table[(v : β | τ | α | κ | ε | η)i → vnei, {i, {is}}]
  ] // LZipJoin@Table[{βnei, τnei, anei}, {i, {is}}] //  

  QZipJoin@Table[{εnei, γnei, {i, {is}}] ];
B{is_}[L_, R_] := B{is}[L, R];

```

“**E** morphisms with domain and range.

```

Bis_List[Ed1 → r1_[L1_, Q1_, P1_], Ed2 → r2_[L2_, Q2_, P2_]] :=

E(d1 ∪ Complement[d2, is]) → (r2 ∪ Complement[r1, is]) @@

Bis[E[L1, Q1, P1], E[L2, Q2, P2]];

Ed1 → r1_[L1_, Q1_, P1_] // Ed2 → r2_[L2_, Q2_, P2_] :=

Br1 ⊗ d2 [Ed1 → r1_[L1, Q1, P1], Ed2 → r2_[L2, Q2, P2]];

Ed1 → r1_[L1_, Q1_, P1_] ≡ Ed2 → r2_[L2_, Q2_, P2_] ^:=

(d1 == d2) ∧ (r1 == r2) ∧ (E[L1, Q1, P1] ≡ E[L2, Q2, P2]));

Ed1 → r1_[L1_, Q1_, P1_] Ed2 → r2_[L2_, Q2_, P2_] ^:=

E(d1 ∪ d2) → (r1 ∪ r2) @@ (E[L1, Q1, P1] × E[L2, Q2, P2]));

Edr_[L_, Q_, P_]$k := Edr @@ E[L, Q, P]$k;

E[E—][i_] := {E}[[i]];

```

E[**Λ**]

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Edr_[A_] :=

CF@Module[{L, Δθ = Limit[A, ε → 0]},
  Edr [L = Δθ /. (η | y | ξ | x) → 0, Δθ - L, eA - Δθ] $k /.12U]

```

Exponentials as needed.

Task. Define $\text{Exp}_{m,i,k}[P]$ to compute $e^{\Omega(P)}$ to ϵ^k in the using the $m_{i,i \rightarrow i}$ multiplication, where P is an ϵ -dependent near-docile element, giving the answer in \mathbb{E} -form.

Methodology. If $P_0 := P_{\epsilon=0}$ and $e^{\lambda \Omega(P)} = \mathcal{O}(e^{\lambda P_0} F(\lambda))$, then

$F(\lambda = 0) = 1$ and we have:

$\mathcal{O}(e^{\lambda P_0}(P_0 F(\lambda) + \partial_\lambda F)) = \mathcal{O}(\partial_\lambda e^{\lambda P_0} F(\lambda)) =$
 $\partial_\lambda \mathcal{O}(e^{\lambda P_0} F(\lambda)) = \partial_\lambda e^{\lambda P_0} = e^{\lambda P_0} \mathcal{O}(P) = \mathcal{O}(e^{\lambda P_0} F(\lambda)) \mathcal{O}(P)$

This is a linear ODE for F . Setting inductively $F_k = F_{k-1} + \epsilon^k \varphi$ we find that $F_0 = 1$ and solve for φ .

```

(* Bug: The first line is valid only if O(eP0) == eO(P0). *)
Expm,i,0[P_] := Module[{LQ = Normal@P /. ε → 0},
  E[LQ /. (x | y)i → 0, LQ /. (b | a | t)i → 0, 1]];
Expm,i,k[P_] := Block[{$k = k},
  Module[{P0, λ, φ, qs, F, j, rhs, eqn, pows, at0, atλ},
    P0 = Normal@P /. ε → 0;
    F = Normal@Last@Expm,i,k-1[λ P];
    While[
      rhs =
        mi,j→i[E{}→{i}[λ P0 /. (x | y)i → 0, λ P0 /. (b | a | t)i → 0,
          Fk sσi,j@E{}→{i}[0, 0, Pk] // Last // Normal];
        eqn = CF[∂λ F + P0 F - rhs];
        eqn != ! 0, (*do*)
        pows = First /@ CoefficientRules[eqn, {yi, bi, ai, xi}];
        F += Sum[ek φjs[λ] Times @@ {yi, bi, ai, xi}js,
          {js, pows}];
        rhs =
        mi,j→i[E{}→{i}[λ P0 /. (x | y)i → 0, λ P0 /. (b | a | t)i → 0,
          Fk sσi,j@E{}→{i}[0, 0, Pk] // Last // Normal];
        eqn = CF[∂λ F + P0 F - rhs];
        qs = Table[φjs[λ], {js, pows}];
        at0 = Table[φjs[0] = 0, {js, pows}];
        atλ = (# == 0) & /@
          (pows /.CoefficientRules[eqn, {yi, bi, ai, xi}]);
        F = F /.DSolve[And @@ (at0 ∪ atλ), φs, λ] [[1]];
      ];
      E{}→{t}[P0 /. (x | y)i → 0, P0 /. (b | a | t)i → 0,
        F + O[ek+1 /.λ → 1]] ]

```

“Define” Code

Video and more: <http://www.math.toronto.edu/~drorbn/Talks/CRM-1907>,
<http://www.math.toronto.edu/~drorbn/Talks/UCLA-191101>.