

Subspaces & Linear independence

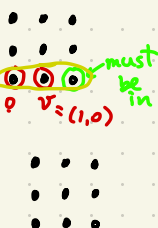
$\mathbb{F}_3^2$ v. space over field $\mathbb{F}_3 = \{0, 1, 2\}$

$$\mathbb{F}_3^2 = \{(a, b) \in \mathbb{F}_3 \times \mathbb{F}_3\}.$$

minimal
 $\{0\} \subseteq V$
subspace



$\text{span}(v)$
 $\{0, v, 2v\}$



$$\begin{aligned} \frac{1}{2} \cdot (2, 2) &= 2 \cdot (2, 2) \\ &= (4, 4) = (1, 1) \end{aligned}$$

$$\text{span}((2, 2)) = \text{span}((1, 1))$$

how many subspaces are there of this type?

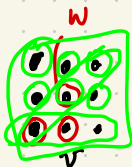
① choose a nonzero vector v
in $\mathbb{F}_p^2$ there are
 $p^2 - 1$ of these

② take scalar multiples of it
 $\Rightarrow$ besides zero there are
 $p - 1$ distinct multiples
of v .

the # of these subspaces is $\frac{p^2 - 1}{p - 1}$ # of nonzero vectors

$$\begin{aligned} &= p + 1 \\ &= 3 + 1 = 4 \end{aligned}$$

of nonzero vectors in $\text{span}(v)$ $v \neq 0$.

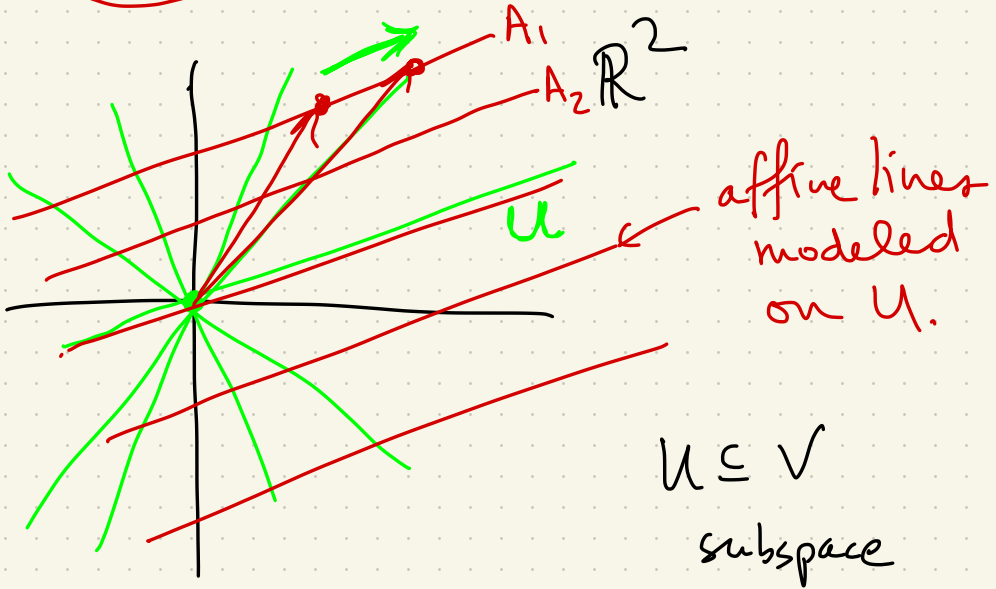
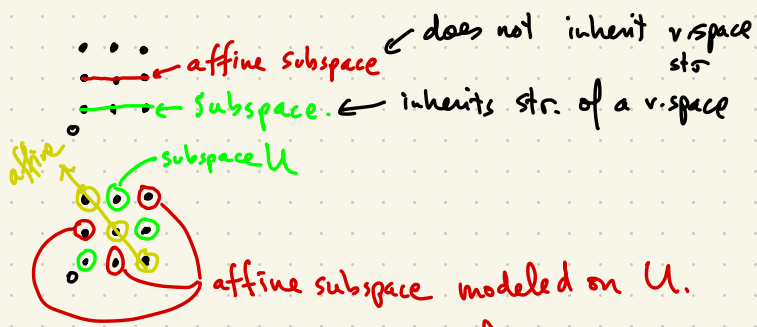


choose $w \notin \text{span}(v)$.

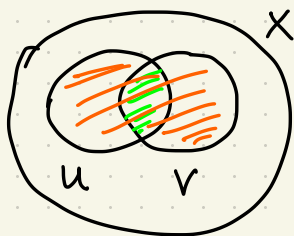
$$\text{span}(v \neq 0, w \notin \text{span}(v)) = \mathbb{F}_3^2$$

check.

$$1 + 1 + 4 = 6$$



Operations on subspaces:



U, V subsets of X

generate new subsets:

- intersection $U \cap V$

"and"

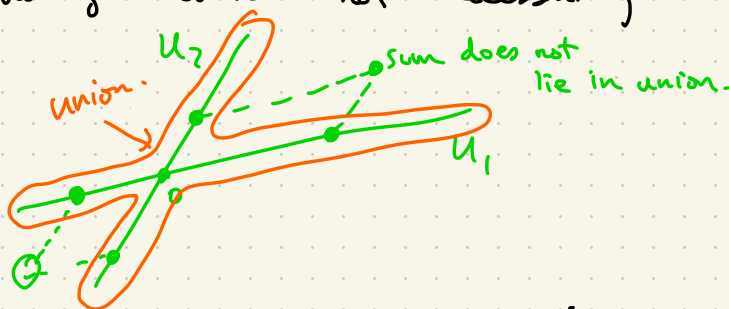
- union $U \cup V$

"or"

Analogy for vector subspaces: if $U_1, U_2 \subseteq V$ are subspaces

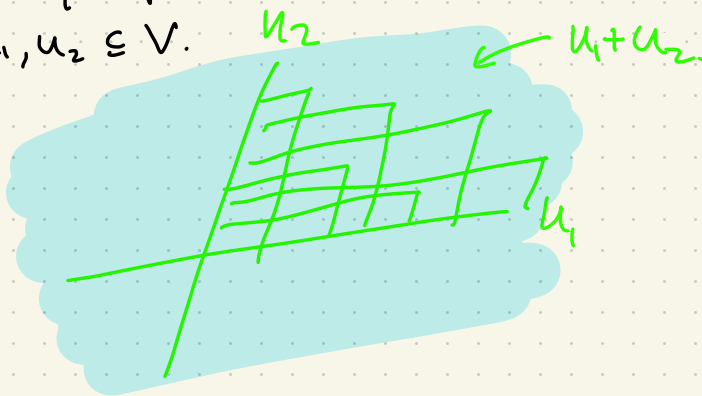
have ① intersection $U_1 \cap U_2 \subseteq V$ is a subspace

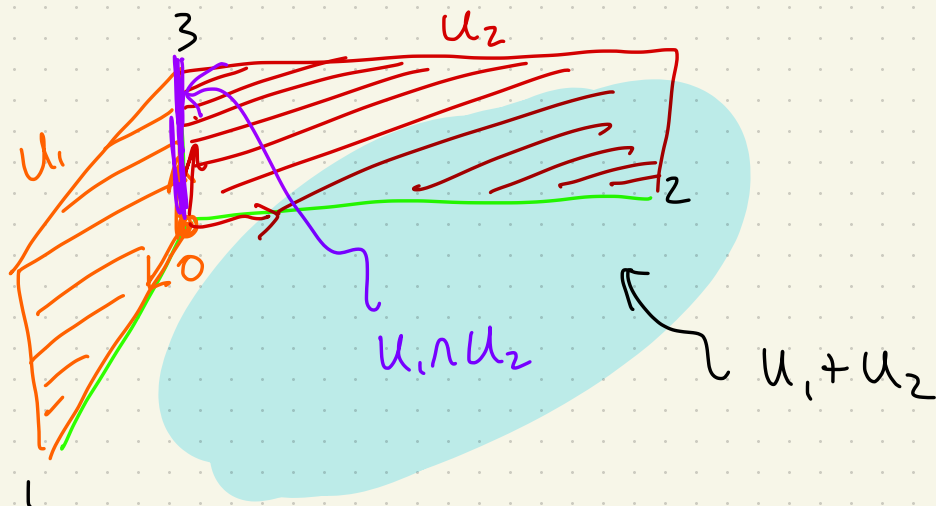
② Warning: union not necessarily a subspace



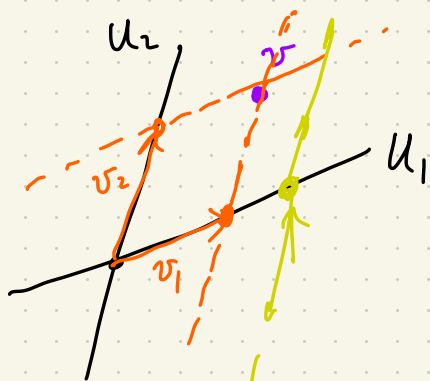
②' Sum of subspaces $U_1 + U_2 = \{ u_1 + u_2 : \begin{matrix} u_1 \in U_1 \\ u_2 \in U_2 \end{matrix} \}$

$u_1, u_2 \in V.$





Property that a sum $U_1 + U_2$ may or may not have: being DIRECT.



$$v_1 + v_2 = v.$$

linear comb, $v_1 \in U_1$
 $v_2 \in U_2$.

Defⁿ

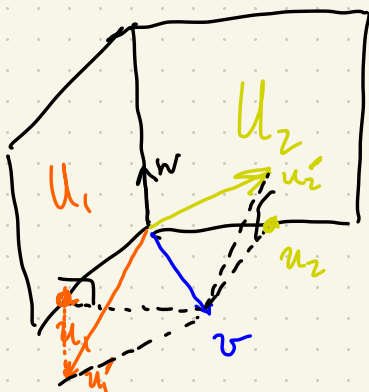
$U_1 + U_2$ is DIRECT (denoted by $U_1 \oplus U_2$)

when any $v \in U_1 + U_2$ has unique expression

$$v = u_1 + u_2$$

$$u_1 \in U_1$$

$$u_2 \in U_2.$$



$$\begin{aligned} u_1 + u_2 &= v \\ u'_1 + u'_2 &= v \end{aligned}$$

where $u'_1 = u_1 + w$
 $u'_2 = u_2 - w$
 $w \in U_1 \cap U_2$

decomp. of
 v
 NOT unique
 so $U_1 + U_2$
 NOT
 DIRECT

$$U_1 + U_2 = \mathbb{R}^3$$

$(v_1, v_2, \dots, v_n)$

Span takes list of vectors in V

outputs a linear subspace $\text{Span}(v_1, \dots, v_n)$
 of V .

Defn

$$\text{Span}(v_1, \dots, v_n) = \{ a_1 v_1 + \dots + a_n v_n : a_i \in \mathbb{F} \}$$

$$\text{Span}(\uparrow) = \{0\}$$

empty list

Note: $\text{Span}(v_1, \dots, v_n) = \text{Span}(v_1) + \text{Span}(v_2) + \dots + \text{Span}(v_n)$

Defⁿ

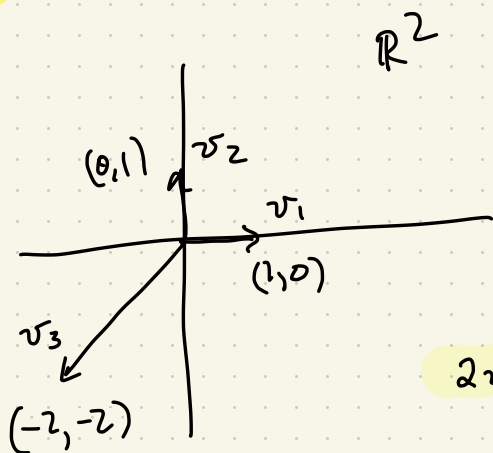
a list $(v_1, \dots, v_n)$ is linearly dependent ↗ "list is redundant"

when there exists $a_1, \dots, a_n$, not all zero,

with
$$a_1 v_1 + a_2 v_2 + \dots + a_n v_n = 0$$

("a nontrivial linear combination equals zero")

If a list is not linearly dep. we say it is
linearly INDEPENDENT.



(v_1, v_2, v_3) is

linearly dependent.

$$a_1 = 2 \quad a_2 = 2 \quad a_3 = 1$$

$$2v_1 + 2v_2 + 1 \cdot v_3 = 0$$

Note: can interpret this as saying $v_3 = (-2)v_1 + (-2)v_2$
i.e. $v_3 \in \text{Span}(v_1, v_2)$

$$\text{or } v_1 = \frac{1}{2}((-2)v_2 + (-1)v_3) = -v_2 - \frac{1}{2}v_3$$

$$\text{i.e. } v_1 \in \text{Span}(v_2, v_3).$$

Note: A list $(v_1, \dots, v_n)$ of vectors in V is linearly independent

$\Leftrightarrow \text{Span}(v_1) + \text{Span}(v_2) + \dots + \text{Span}(v_n)$ is DIRECT.

Defⁿ: V is finite-dimensional when it can be spanned by a finite list of elements

i.e. $V = \text{Span}(v_1, \dots, v_n)$.

(V infinite-dimensional if not).

e.g. $\mathbb{R}^n$ is finite-dimensional

because

$$\mathbb{R}^n = \text{Span}(e_1, \dots, e_n, e_1, e_3, e_2+e_4)$$

$$e_1 = (1, 0, \dots, 0)$$

$$e_2 = (0, 1, 0, \dots, 0)$$

$\vdots$

$$e_n = (0, 0, \dots, 0, 1)$$

e.g. $\mathcal{P}(\mathbb{R})$ polynomials
not finite dim.

(2.21)

Lemma: If $(v_1, \dots, v_n)$ is linearly dependent,

then ① $\exists v_j$ in the span of the previous vectors on list,

i.e. $v_j \in \text{Span}(v_1, \dots, v_{j-1})$

② and we may remove v_j without changing span of list.

Proof: $(v_1, \dots, v_n)$ is lin. dep means $\exists (a_1, \dots, a_n) \in \mathbb{F}^n \setminus \mathbf{0}$

s.t. $a_1 v_1 + \dots + a_n v_n = \mathbf{0}$

(If $\mathbf{0}$ is on the list the list is lin. dep.)

Scan from end $\left\{ a_n, a_{n-1}, a_{n-2}, \dots \right\}$ until we find $a_j \neq 0$

i.e. this eqn is actually $a_1 v_1 + \dots + \overset{\text{nonzero number}}{a_j} v_j = \mathbf{0}$

since $a_j \neq 0$ we can solve for v_j :

$$v_j = -\underset{\substack{\text{exists since } \mathbb{F} \text{ field.}}}{a_j^{-1}} (a_1 v_1 + \dots + a_{j-1} v_{j-1}).$$

$\Rightarrow v_j \in \text{Span}(v_1, \dots, v_{j-1}) \Rightarrow \text{prove } ①$

②: $\text{Span}(v_1, \dots, v_{j-1}, v_{j+1}, \dots, v_n) = A$

claim $\rightarrow \text{Span}(v_1, \dots, v_n) = B$

$A \subseteq B$ obvious

$B \subseteq A$: $w \in B$

$$B \ni w = b_1 v_1 + \dots + b_n v_n$$

$$= b_1 v_1 + \dots + b_{j-1} v_{j-1}$$

$$+ b_j (-a_j^{-1} (a_1 v_1 + \dots + a_{j-1} v_{j-1}))$$

$$+ b_{j+1} v_{j+1} + \dots + b_n v_n.$$

$$\in A.$$

(2.23)

Thm. $\text{len}(\text{lin indep list}) \leq \text{len}(\text{spanning list}).$

i.e. If $V = \text{span}(w_1, \dots, w_n)$

and if $(u_1, \dots, u_m)$ is lin indep.

then $m \leq n.$

Proof
algorithm: use ordering + lemma:

① adjoin u_1 to spanning list:

$(u_1, w_1, \dots, w_n)$ $\begin{cases} \text{still spans } V \\ \text{linearly dependent} \end{cases}$
since $u_1 \in V = \text{Span}(w_1, \dots, w_n)$

by the lemma $\exists w_j$ which is in $\text{span}(u_1, w_1, \dots, w_{j-1})$.

eliminate w_j , the result $(u_1, w_1, \dots, w_{j-1}, w_{j+1}, \dots, w_n)$

still spans V .

Continue this process: $(u_1, u_2, w_1, \dots, \overset{w_j \text{ not in list}}{\hat{w}_j}, \dots, w_n)$ $\begin{cases} \text{still spans } V \\ \text{lin. dependent} \end{cases} \Rightarrow \exists w_k \text{ to eliminate}$
each time add $u \xRightarrow{\text{lemma}} \exists$ vector on list linearly dependent on previous elements
this cannot be one of the u 's by assumption. $\Rightarrow$ must be one of the $w_i \Rightarrow$ remove it.

so for each new u_i , there is a w_j which we can remove
and continue. $\Rightarrow \# u_i \leq \# w_i.$

Review: $(v_1, \dots, v_n)$ lin. dependent i.e. $\exists$ redundancy from the p.o.v of span.

there is some nontrivial lin. comb. $= 0$.

i.e. can find $(a_1, \dots, a_n)$ not all zero st. $a_1 v_1 + \dots + a_n v_n = 0$.

$a_1 v_1 + a_2 v_2 = 0$ $\implies a_2 = -a_1$
 $2v_1 - v_2 = 0$ $\implies a_2 \neq 0$

Def: Let V be a finite dim. v. space (it has a finite spanning list),
 A basis for V is a list $(v_1, \dots, v_m)$ i.e. $V = \text{Span}(v_1, \dots, v_n)$ for some $v_i \in V$.
 of vectors which

- ① spans V
- ② lin. independent.

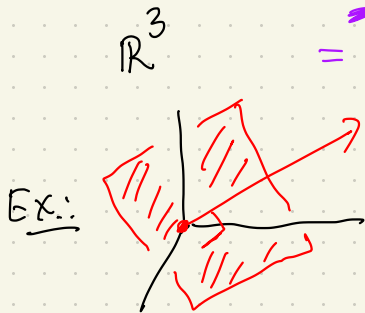
Ex.: $V = \mathbb{F}^n$ has a natural basis as follows:

$$\left. \begin{aligned} e_1 &= (1, 0, 0, \dots, 0) \\ e_2 &= (0, 1, 0, \dots, 0) \\ &\vdots \\ e_n &= (0, 0, \dots, 0, 1) \end{aligned} \right\}$$

① any vector in V can be written as
 $v = (a_1, \dots, a_n)$ $a_i \in \mathbb{F}$
 $v = a_1 e_1 + a_2 e_2 + \dots + a_n e_n$
 i.e. $\text{Span}(e_1, \dots, e_n) = V$.

$$\lambda_1 (1, 0, 0) + \lambda_2 (0, 1, 0) + \lambda_3 (0, 0, 1) = (\lambda_1, \lambda_2, \lambda_3) = (0, 0, 0)$$

② $(e_1, \dots, e_n)$ is lin. indep.
 If $\lambda_1 e_1 + \dots + \lambda_n e_n = 0$ $\lambda_i \in \mathbb{F}$
 $\parallel$
 $(\lambda_1, \dots, \lambda_n) = 0 = (0, 0, \dots, 0)$
 $\implies \lambda_i = 0 \quad \forall i$.
 (linear comb is triv. no redundancy)



Ex.: $U \subset \mathbb{R}^3$ lin subspace $U = \{(x, y, z) : x + y + z = 0\}$

Q.: find a basis for U ?

(Pessimist approach)

Thm: Every finite dimensional v.space has a basis (highly non-unique choice).

Pf: If $V = \{0\}$ then the empty list $()$ is a basis.
otherwise can express $V = \text{Span}(v_1, \dots, v_n)$ v_i nonzero.
for each $v_2, v_3, \dots, v_n$ we ask if it is in the span
of previous vectors — if so, remove it
— if not keep it.

after this, span is unchanged, and no vector is in
span of previous. By Lemma the resulting list is lin. indep.

$\Rightarrow$ lin indep. list spans $V \Rightarrow$ BASIS. $\square$

Thm: (Build basis)
optimist approach Any lin. indep. list in a finite dim. v.sp.
can be extended to a basis

Pf: If $(v_1, \dots, v_m)$ is lin. independent, let $(w_1, \dots, w_n)$
be a spanning list.

Step 1: check if $w_1 \in \text{Span}(v_1, \dots, v_m)$ If so, leave B
unchanged. If not, add w_1 to B .
new $B = (v_1, \dots, v_m, w_1)$.

Step j If w_j in span of the latest B (from step $j-1$)
leave unchanged, if not, add as before.

finally after n steps: • have (by Lemma) a lin indep. list.
list looks like
 $(v_1, \dots, v_m, w_{i_1}, w_{i_2}, \dots, w_{i_k})$
• but $\text{Span}(w_1, \dots, w_n) \subseteq \text{Span}(v_1, \dots, v_m, w_{i_1}, \dots, w_{i_k})$
" V

$\mathbb{R}^3$

$()$

lin. indep

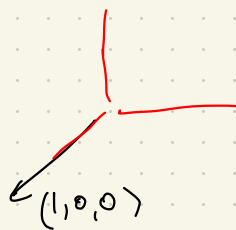
(x, y, z)

$(1, 0, 0)$

lin. indep.

$x, y, z \in \mathbb{R}$

$$a(1, 0, 0) = 0$$



$$a = 0$$

$$\left(\overset{v_1}{(1, 0, 0)}, \overset{v_2}{(1, 1, 0)} \right)$$

$$a(1, 0, 0) + b(1, 1, 0) = (0, 0, 0)$$

$$\parallel$$
$$(a, 0, 0) + (b, b, 0)$$

$$(a+b, b, 0) = (0, 0, 0)$$

$$1) b = 0$$

$$2) a = 0.$$

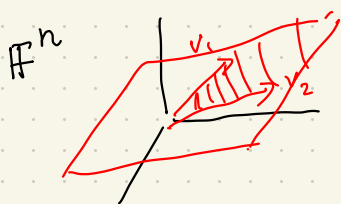
$$\left(\overset{v_1}{(1, 0, 0)}, \overset{v_2}{(1, 1, 0)}, \overset{v_3}{(1, 1, 1)} \right) \underline{\text{Basis}}$$

$$\text{Span} = \mathbb{R}^3$$

Defⁿ: Let V be finite dimensional
 (Incomplete) v.s. space. The dimension of V
 is the length of a basis
 for V .

This definition depends on the fact that all bases^{of V} have
 the same length — otherwise "ill-defined".

Comments about explicit subspaces:



Subspaces: $\text{Span}(v_1, \dots, v_k)$

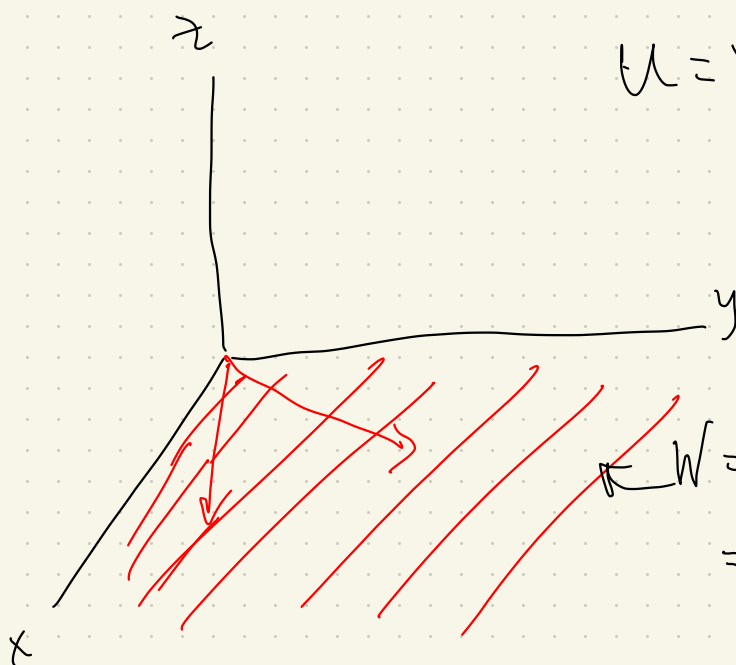
$$\begin{aligned} v_1 &= (a_1, \dots, a_n) & a_i &\in \mathbb{F} \\ v_2 &= (b_1, \dots, b_n) & b_i &\in \mathbb{F} \\ &\vdots & & \\ v_k &= (c_1, \dots, c_n) & c_i &\in \mathbb{F} \end{aligned}$$

$k \times n$ grid of numbers

alternatively: $(a_1, \dots, a_n) \quad a_i \in \mathbb{F}$
 $(b_1, \dots, b_n) \quad b_i \in \mathbb{F}$

$$\left\{ (x_1, \dots, x_n) \in \mathbb{F}^n : \begin{array}{l} \checkmark a_1 x_1 + \dots + a_n x_n = 0 \\ \checkmark b_1 x_1 + \dots + b_n x_n = 0 \\ \checkmark \vdots \end{array} \right\}$$

constraint



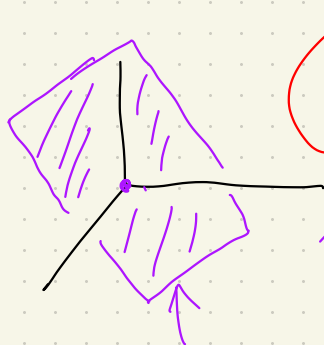
$U = \underline{xy \text{ plane}} \subseteq \mathbb{R}^3$
 subspace.

1. as a span

$$\begin{aligned} W &= \text{Span}((1, 1, 0), (1, 2, 0)) \\ &= \text{Span}((1, 0, 0), (0, 1, 0)) \\ &\rightarrow \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix} \\ &\rightarrow \begin{bmatrix} 1 & 1 & 0 \\ 1 & 2 & 0 \end{bmatrix} \end{aligned}$$

2. as a constraint

$$\{(x, y, z) : 0 \cdot x + 0 \cdot y + 1 \cdot z = 0\} \quad [0 \ 0 \ 1]$$



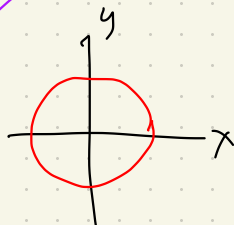
$$U = \{(x, y, z) :$$

Solutions to this = linear subsp.

$$x + y + z = 0\}$$

Linear equations.

~~$$\{(x, y) : x^2 + y^2 = 1\}$$~~



$$(1, 0, 0)$$

$$(0, 1, 0)$$

$$(0, 0, 1)$$

Q: find a basis for U

$$(x, y, z) \in W$$

$\Rightarrow$

$$0 \cdot x + 0 \cdot y + 1 \cdot z = 0$$

and

$$1 \cdot x + 1 \cdot y + 1 \cdot z = 0$$

$$(x, y, z) \in U \Rightarrow$$

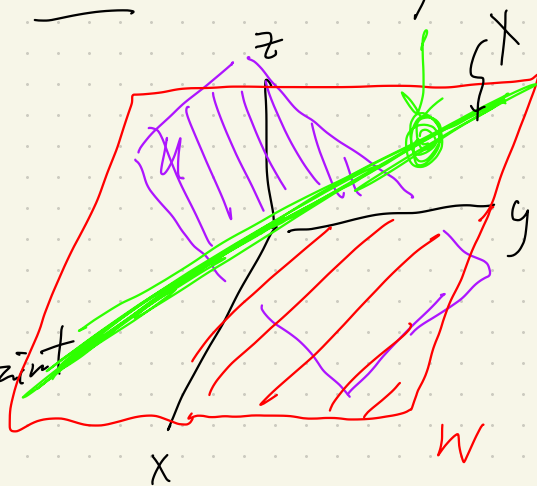
defines a linear subspace of $\mathbb{R}^3$

$$\{(x, y, z) \text{ st. in } W \text{ and in } U\}$$

$$X = U \cap W$$

1st constraint

2nd constraint



$\mathcal{P}_{\leq k}(\mathbb{F})$

polynomials of maximum degree k
w/ coeffs in $\mathbb{F}$.

Q.: what is the dimension?

Basis $(1, x, x^2, \dots, x^k)$

$$\dim(\mathcal{P}_{\leq k}(\mathbb{F})) = k+1$$

$$\left\{ \begin{array}{l} 1 \\ x+1 \\ x^2+x+1 \\ x^3+x^2+x+1 \\ \vdots \\ x^k+x^{k-1}+\dots+x+1 \end{array} \right.$$

$\mathcal{P}(\mathbb{F})$ v.space over $\mathbb{F}$

$(v_1, \dots, v_m)$ $v_i \in \mathcal{P}(\mathbb{F})$. v_i has max degree d_i

st. $\mathcal{P}(\mathbb{F}) = \text{Span}(v_1, \dots, v_m)$.

$$a_1 v_1 + \dots + a_m v_m$$

$$v_1 = a_0 + a_1 x + a_2 x^2 + \dots + a_{d_1} x^{d_1}$$

v_i has a max degree d_i

$$\lambda v_1 = (\lambda a_0) + (\lambda a_1) x + \dots + (\lambda a_{d_1}) x^{d_1}$$

$$\lambda_1 v_1 + \dots + \lambda_n v_n \quad \vec{x} \mapsto \vec{x}, \vec{a}$$

linear constraints of the form

$$a_1 x_1 + \dots + a_n x_n = 0$$

$$\mathbb{F}^n \ni (x_1, \dots, x_n)$$

$$a_i \in \mathbb{F}$$

this constraint is really $f^{-1}(0)$

$$f: \mathbb{F}^n \longrightarrow \mathbb{F}$$

$$(x_1, \dots, x_n) \mapsto a_1 x_1 + \dots + a_n x_n$$

Span $\longrightarrow$ Finite dimensionality.

Linear (in)dependence $\longrightarrow$ basis (lin ind + spanning)

lemma: length (linearly indep list) $\leq$ length (spanning list)

$\Rightarrow$ Defⁿ of Dimension

Def: $\dim V = \text{len}(\text{basis})$ $\swarrow$ V has a basis! (Lemma)
this makes sense only if $\swarrow$ all bases have same length!

lemma: If B, B' are bases for $V \Rightarrow \text{len}(B) = \text{len}(B')$

PF: B spans B' lin indep $\Rightarrow \text{len}(B') \leq \text{len}(B)$
Lemma

B' spans B lin indep $\Rightarrow \text{len}(B) \leq \text{len}(B')$
Lemma

$\Rightarrow \text{len}(B) = \text{len}(B')$

ex: $\mathcal{P}_{\leq n}(\mathbb{F}) = \{ a_0 + a_1x + \dots + a_nx^n : a_i \in \mathbb{F} \}$

has a basis $B = (p_0, p_1, \dots, p_n)$

$p_0 = 1$
 $p_1 = x$
 $p_2 = x^2$
 $\vdots$
 $p_n = x^n$

Basis

Spans:

$$= a_0p_0 + a_1p_1 + \dots + a_np_n$$

lin. indep:

$$a_0 + a_1x + \dots + a_nx^n = 0$$

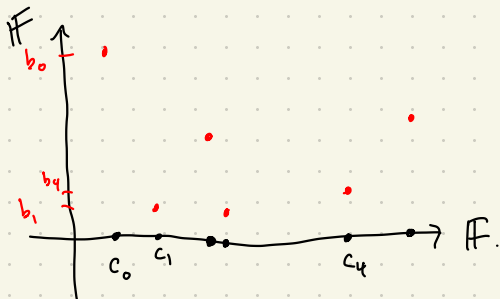
$$\Leftrightarrow a_0 = a_1 = a_2 = \dots = a_n = 0$$

$$\text{len } B = n+1$$

$$\dim \mathcal{P}_{\leq n}(\mathbb{F}) = n+1$$

Ex.: $\mathcal{P}_n(\mathbb{F})$ same vector space, different basis.

$$\dim \mathcal{P}_n(\mathbb{F}) = n+1.$$



Problem: given input numbers
 $c_0, c_1, \dots, c_n \in \mathbb{F}$
 and output values.

$$b_0, \dots, b_n,$$

Can we find a polynomial
 which fits the data.

i.e. p s.t.

$$p(c_i) = b_i \quad ?$$

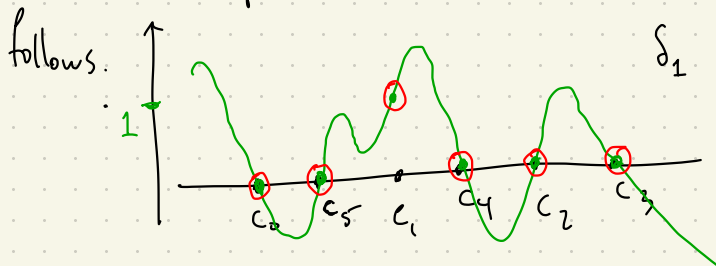
heuristic:

Guess: this might be possible with $\mathcal{P}_n(\mathbb{F})$
 since $\dim = n+1$.

Reason why problem is hard: basis we are used to is
 $(1, x, x^2, \dots, x^n)$

so unclear how choice of coeffs $a_0 + a_1x + \dots + a_nx^n$
 affects value of poly at $c_0, \dots, c_n$.

Alternatively: Suppose we had a different
 basis $(\delta_0, \dots, \delta_n)$ adapted to positions $c_0, \dots, c_n$ as



$$\delta_1(c_0) = 0$$

$$\delta_1(c_1) = 1$$

$$\delta_1(c_2) = 0$$

$\vdots$

$$\delta_1(c_n) = 0.$$

have "indicator basis" adapted to $(c_0, \dots, c_n)$

$$\delta_i(c_j) = \begin{cases} 1 & \text{if } i=j \\ 0 & \text{else.} \end{cases}$$

Can we solve using this basis?

need p st. $p(c_i) = b_i$

$b_i \delta_i$ has values

0	at	c_0
0	at	c_1
0	at	c_{i-1}
b_i	at	c_i
0	at	c_{i+1}
$\vdots$		
0	at	c_n

conclude $b_0 \delta_0 + \dots + b_n \delta_n = p$!!

↑ satisfies all constraints!

how do we find such a basis?

(Lagrange Interpolation)
data fitting.

δ_i should have value 1 at c_i and 0 at c_j $j \neq i$

$$\delta_i = \frac{(x - c_0)(x - c_1) \dots (x - c_{i-1})(x - c_{i+1}) \dots (x - c_n)}{(c_i - c_0)(c_i - c_1) \dots (c_i - c_{i-1})(c_i - c_{i+1}) \dots (c_i - c_n)}$$

has zeros
at all
 c_j $j \neq i$.

at c_i it has value

this will work if the c_i are pairwise distinct $(c_i - c_0)(c_i - c_1) \dots$

Summary: $\left\{ \begin{array}{l} \text{Fix pairwise distinct elts of } \mathbb{F} \\ c = (c_0, \dots, c_n) \\ \text{Create basis } B_c = (\delta_0, \dots, \delta_n) \\ \text{s.t. } \delta_i(c_j) = \begin{cases} 1 & i=j \\ 0 & \text{else.} \end{cases} \end{array} \right.$

$$\delta_i = \prod_{\substack{j=0 \\ j \neq i}}^n \frac{(x - c_j)}{(c_i - c_j)}$$

Fit to data is $\sum_{i=0}^n b_i \delta_i = p. \in \mathcal{P}_n(\mathbb{F})$

$$p(c_i) = b_i$$

Rem: How do we know it's a basis? It is lin indep.
Its length

It is linearly independent!:

is $n+1$
 $\downarrow$
basis

If $(a_0, \dots, a_n)$ are such that

$$f = a_0 \delta_0 + \dots + a_n \delta_n \stackrel{(*)}{=} 0$$

want to show all $a_i = 0 \Rightarrow$ list B_c lin indep.

but $\Rightarrow \boxed{f(c_i) \stackrel{\forall i}{=} a_i = 0}$

$\square$